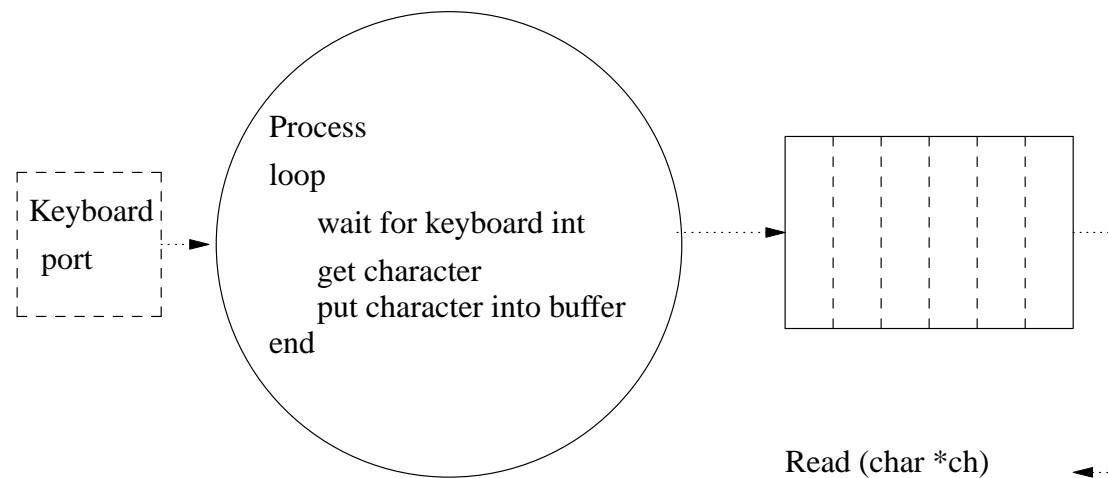


Shared buffer

- laboratory 2 implements a shared buffer



- shared by the application process `Read (char *ch) ;`
 - shared by the driver process
-
- what are the problems with this?
 - data might be altered by both processes at the same time
 - data (buffer pointers) could become corrupt

Mutual exclusion

- require a mechanism to ensure that only one process can manipulate data at any one time
 - *mutual exclusion*
- the concepts we discuss today are *very* important for microkernels and operating systems
 - a fundamental building block

How do we implement mutual exclusion?

- simplest mechanism
 - mask processor interrupts off
 - processor cannot respond to any interrupt and therefore will execute code in sequence until it masks interrupt back on again
 - sometimes these critical sections of code are called *atomic*
 - what are this disadvantages with this approach?
 - what are this advantages with this approach?

How do we implement mutual exclusion?

- another mechanism is *semaphores*
 - essentially a binary *semaphore* is a token which can be grabbed by *only one* process at a time
 - a token is taken at the entry to the critical section and given back at the end of the critical section
 - a process can only enter once it has the token

Semaphores

- consider the following two processes:

```
        (* Shared semaphore *)
Semaphore token;  (* initial value 1 *)

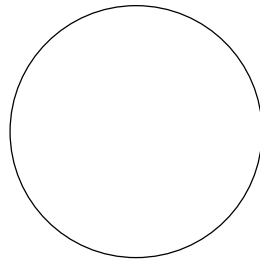
void ProcessA ()          void ProcessB ()
{                          {
    while (TRUE) {        while (TRUE) {
        ...                ...
        Wait(Token)        Wait(Token)
        (* critical *)    (* critical *)

        Signal(Token)     Signal(Token)
        ...                ...
    }                      }
}                          }
```

Semaphores



SEMAPHORE token



- Wait gets the token
- Signal returns the token

Semaphores

- note that `Wait` and `Signal` are both *atomic*
- they are implemented in software with processor interrupts masked off

Semaphores

- we can express Wait and Signal in pseudo code:

```
void Wait (s)
{
    when s>0
        s--;
}

void Signal (s)
{
    s++;
}
```

Semaphores

- in our previous example the initial value of s would be 1
 - note that this is pseudo code
 - note the use of **when**

Semaphores

- we have now seen how a critical section can be achieved by using semaphore primitives `Wait` and `Signal`
- for example access to the shared buffer will be a critical section

Starting to implement a shared buffer using semaphores

```
void put (char ch)          char get (void)
{
    Wait (Mutex)           Wait (Mutex)
    (* safe to alter *)   (* safe to alter *)
    (* buffer      *)     (* buffer      *)
    place ch into buf     remove ch from buf

    Signal (Mutex)        Signal (Mutex)

                                return ch;
}

char buffer[Max]; (* global data *)
SEMAPHORE Mutex; (* global data *)
```

Completed implementation of a shared buffer using semaphores

```
void put (char ch)          char get (void)
{
    Wait (SpaceAvailable)   {
    Wait (Mutex)            Wait (ItemAvailable)
                            Wait (Mutex)

    (* safe to alter *)    (* safe to alter *)
    (* buffer      *)      (* buffer      *)
    place ch into buf      remove ch from buf

    Signal (Mutex)         Signal (Mutex)
    Signal (ItemAvailable) Signal (SpaceAvailable)
                            return ch;
}

char buffer[Max]; (* global data *)
SEMAPHORE Mutex; (* global data *)
```

Shared buffers

- see `lab/boundedbuffer/BufferDevice.c` for a shared buffer data structure
- remember we need to solve the problem of what should happen if:
 - there is no data to take out?
 - there is no room left in the buffer?

Shared buffers

- if there is no data in the buffer and we attempt to get a datum then we should *wait* until data arrives
- if there is no space in the buffer and we attempt to put a datum then we should *wait* until space available
- we can implement this with two semaphores
 - ItemAvailable
 - SpaceAvailable

Shared buffer (continued)

- before we place an item into a buffer we must
`Wait (SpaceAvailable)`
- before we extract an item from a buffer we must
`Wait (ItemAvailable)`
- after we place an item into the buffer we must
`Signal (ItemAvailable)`
- after we extract an item from the buffer we must
`Signal (SpaceAvailable)`

Shared buffer (continued)

- what are the initial values for an empty buffer? (size 3)
 - `ItemAvailable` 0
 - `SpaceAvailable` 3

- this buffer mechanism is known as Dijkstra's bounded buffer after its author E.W. Dijkstra who discovered the algorithm in 1960s

Shared buffer (continued)

Wait (SpaceAvailable) Wait (ItemAvailable)

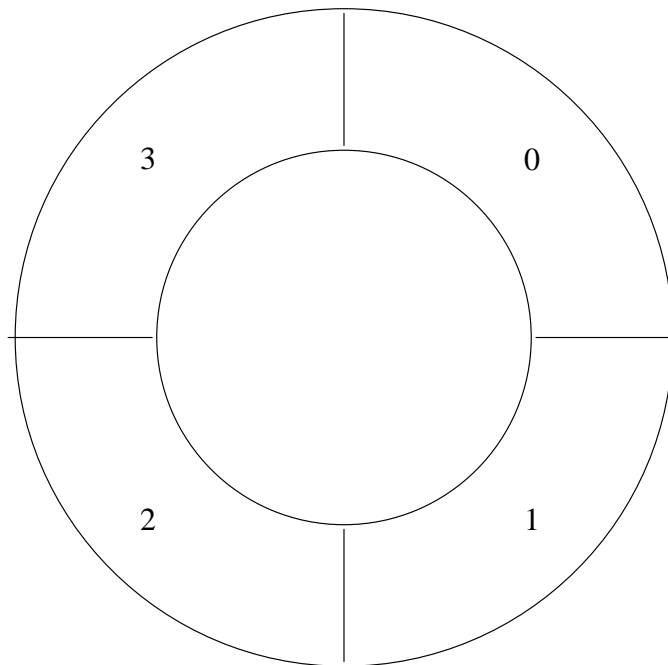
Signal (ItemAvailable) Signal (SpaceAvailable)

Semaphores

- the module `Executive.c` in the `luk` documentation exports the:
 - type `SEMAPHORE`
 - functions `Wait` and `Signal`
- you can use these functions to implement your bounded buffer

Circular buffer details

- circular buffer manipulation
 - need to put a datum in at the front
 - need to extract a datum from the end



Circular buffer details

- implement this with two indices
 - `in` and `out`
 - when we add a datum to the buffer we place it at position `in`. We then increment `in` modulo the buffer size.
 - when we extract a datum from the buffer we extract it from position `out`. We then increment `out` modulo the buffer size.

Circular Buffer Code

- when we add a datum to the circular buffer we:

```
Buf[in] = ch;  
in = (in+1) % MaxBufferSize;
```

- when we extract a datum from a circular buffer we:

```
ch = Buf[out];  
out = (out+1) % MaxBufferSize;
```

Circular Buffer Code (continued)

- in this laboratory tutorial you have write the procedure `InitBuffer`. It must perform 3 operations:
 - it must create a buffer
 - initialize the buffer pointers to correct values
 - initialize the semaphore values

- create the buffer.

Creating and initializing a buffer

- the data structure `Buffer` is a pointer type
 - you must make sure it points to something sensible!
 - to do this use the function `Storage_ALLOCATE`

```
Storage_ALLOCATE((void **) &b, sizeof(Buffer));
```

- initialize the buffer pointers `in` and `out`

```
b->in = 0;  
b->out = 0;
```


Circular Buffer Code (continued)

- initialize semaphore values
 - *must use procedure* `InitSemaphore` from `Executive.c` to initialize semaphores!

- example

- ```
b->Mutex = Executive_InitSemaphore(1, SafeStr("Mutex"));
```
  
- you must initialize the two other semaphores
  - `ItemAvailable`
  - `SpaceAvailable`
  - to their respective values 0 and `MaxBufferSize`

## Device driver (revisited)

- recall that the high level description of the device driver was:

```
void DeviceDriver (void)
{
 (* mask processor *)
 (* ints off *)
 TurnInterruptsOff ;

 SetupDevice ;
 EnableDeviceInterrupts ;
 while (TRUE) {
 WaitForInterrupt (DeviceInterrupt) ;
 ServiceDevice ;
 Store or retrieve data;
 }
}
```

- the value of DeviceInterrupt is 021H for our microkernel

## Laboratory: bounded buffer

- once you have completed this exercise you should be able to type characters and see them appear on the screen
  
- two processes active on your microkernel
  - device driver `ReadDriver`
  - user code (reads `ch` and writes `ch`)
  
- remember that your device driver is responding to interrupts and successfully placing characters into the shared buffer
  
- the application process will extract characters from the buffer and display them to the screen

## Self check work and extra work

- describe major data structures being used by: `BufferDevice.c`,  
`Scn.c`
- make a note of each module - what is its primary purpose
- if you have more time, provide a commentary on the boot code