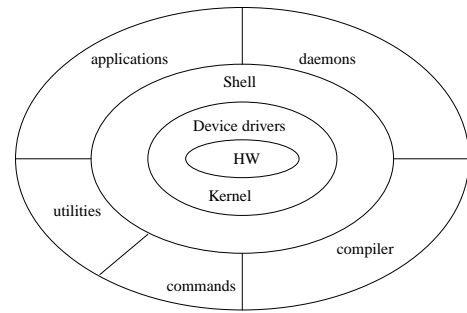


Interrupt handling and context switching

- these two topics are separate and we will examine them in turn

Interrupts

■



- the user programs and hardware communicates with the kernel through interrupts

Four different kinds of interrupts

- device interrupt, such as a hardware timer, for example the 8253 counter0 reaching 0 on an IBM-PC
- user code issuing a software interrupt, often called a **system call**
- an illegal instruction (divide by zero, or an opcode which the processor does not recognise)
- or a memory management fault interrupt (occurs when code attempts to read from non existent memory)

First level interrupt handler

- the kernel must detect which kind of interrupt has occurred and call the appropriate routine
 - this code is often termed the **first level interrupt handler**
- the pseudo code for the FLIH follows:

First level interrupt handler

- ```
save program registers and disable interrupts
k = get_interrupt_kind ();
if (k == source 1) service_source1 ();
else if (k == source 2) service_source2 ();
else if (k == source 3) service_source3 ();
else if (k == source 4) service_source4 ();
else if (k == source 5) service_source5 ();
 etc
restore program registers and enable interrupts
return
```
- you may find the hardware on the microprocessor performs the save and restore program registers and disabling/enabling interrupts
  - possibly by one instruction

## First level interrupt handler

- you might also find the hardware enables you to determine the source of the interrupt easily
  - most microprocessors have an interrupt vector table
    - typically one vector per source is implemented
- equally, however the code can be ugly as it depends upon the hardware specifications

## Example of FLIH in GNU LuK

- GNU LuK (Lean uKernel) is a very small microkernel which allows preemptive processes, interrupt driven devices and semaphores

## Example of FLIH in GNU LuK

- ```
IsrTemplate[ 0] := 0FCH ; (* cld (disable interrupts)
IsrTemplate[ 1] := 050H ; (* push eax *)
IsrTemplate[ 2] := 051H ; (* push ecx *)
IsrTemplate[ 3] := 052H ; (* push edx *)
IsrTemplate[ 4] := 01EH ; (* push ds *)
IsrTemplate[ 5] := 006H ; (* push es *)
IsrTemplate[ 6] := 00FH ; (* push fs *)
IsrTemplate[ 7] := 0A0H ;
IsrTemplate[ 8] := 0B8H ; (* movl 0x00000010, %eax *)
IsrTemplate[ 9] := 010H ;
IsrTemplate[10] := 000H ;
IsrTemplate[11] := 000H ;
IsrTemplate[12] := 000H ;
IsrTemplate[13] := 08EH ; (* mov ax, ds *)
IsrTemplate[14] := 0D8H ;
IsrTemplate[15] := 08EH ; (* mov ax, es *)
IsrTemplate[16] := 0C0H ;
IsrTemplate[17] := 08EH ; (* mov ax, fs *)
IsrTemplate[18] := 0E0H ;
```

Example of FLIH in GNU LuK

```

IsrTemplate[19] := 068H ; (* push interruptnumber *)
IsrTemplate[20] := 000H ; (* vector number to be overwriten *)
IsrTemplate[21] := 000H ; (* this is the single parameter *)
IsrTemplate[22] := 000H ; (* to function. *)
IsrTemplate[23] := 000H ;
IsrTemplate[24] := 0B8H ; (* movl function, %eax *)
IsrTemplate[25] := 000H ; (* function address to be overwritten *)
IsrTemplate[26] := 000H ;
IsrTemplate[27] := 000H ;
IsrTemplate[28] := 000H ;

```

Example of FLIH in GNU LuK

```

IsrTemplate[29] := 0FFH ; (* call %eax *)
IsrTemplate[30] := 0D0H ;
IsrTemplate[31] := 058H ; (* pop %eax // remove parameter *)
IsrTemplate[32] := 00FH ; (* pop %fs *)
IsrTemplate[33] := 0A1H ;
IsrTemplate[34] := 007H ; (* pop %es *)
IsrTemplate[35] := 01FH ; (* pop %ds *)
IsrTemplate[36] := 05AH ; (* pop %dx *)
IsrTemplate[37] := 059H ; (* pop %cx *)
IsrTemplate[38] := 058H ; (* pop %ax *)
IsrTemplate[39] := 0CFH ; (* iret *)

```

Example of FLIH in GNU LuK

- GNU LuK uses a routine `ClaimIsr` which will copy the `IsrTemplate` into the correct interrupt vector and then overwrite the vector number and function address in the template

Context switching

- the scheduler runs inside the kernel and it decides which process to run at any time
 - processes might be blocked waiting on a semaphore or waiting for a device to respond
 - a process might need to be preemptively interrupted by the scheduler if it were implementing a round robin algorithm
- the minimal primitives to manage context switching in a microkernel or operating system were devised by Wirth 1983 (Programming in Modula-2)
 - `NEWPROCESS`, `TRANSFER` and `IOTRANSFER` (covered later on)

A tiny example of two simple processes in an operating system

```

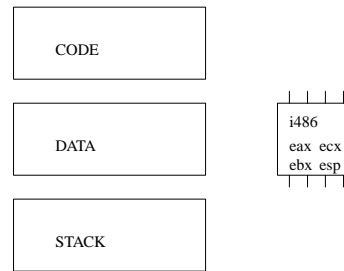
void Process1 (void)
{
    while (TRUE) {
        WaitForACharacter();
        PutCharacterIntoBuffer();
    }
}

void Process2 (void)
{
    while (TRUE) {
        WaitForInterrupt();
        ServiceDevice();
    }
}

```

Primitives to manage context switching

- firstly let us look at a conventional program running in memory (single program running on a computer)



Primitives to manage context switching

- four main components
 - code
 - data
 - stack
 - processor registers (volatiles)

Concurrency

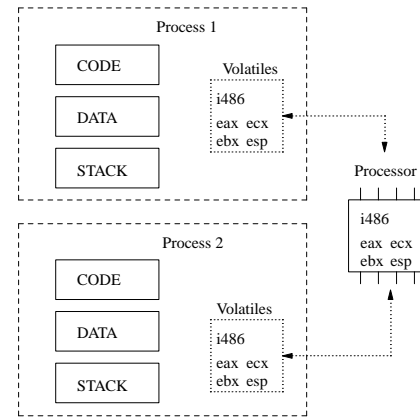
- suppose we want to run two programs concurrently?
 - we could have two programs in memory. (Two stacks, code, data and two copies of a volatile environment)
 - on a single processor computer we can achieve apparent concurrency by running a fraction of the first program and then run a fraction of the second.
 - if we repeat this then apparent concurrency will be achieved
 - in operating systems multiple concurrent programs are often called *processes*

Concurrency

- what technical problems need to be solved so achieve apparent concurrency?
 - require a mechanism to switch from one process to another
- remember our computer has one processor but needs to run multiple processes
 - the information about a process is contained within the volatiles (or simply: processor registers)

Implementing concurrency

- we can switch from one process 1 to process 2 by:
 - copying the current volatiles from the processor into an area of memory dedicated to process 1
 - now copying some new volatiles from memory dedicated to process 2 into the processor registers



Implementing concurrency

- this operation is call a context switch (as the processors context is switched from process 1 to process 2)
 - by context switching we have a completely new set of register values inside the processor
 - so on the i486 we would change **all** the registers. Some of which include: EAX, EBX, ECX, EDX, ESP and flags
 - note that by changing the ESP register (stack pointer) we have effectively changed stack

Context switching primitives in GNU LuK

- the previous description of context switching is very low level
- in a high level language it is desirable to avoid the assembler language details as far as possible
 - NEWPROCESS
 - TRANSFER
 - IOTRANSFER
- **it is possible to build a microkernel which implements context switching and interrupt driven devices using these primitives without having to descend into assembly language**
 - these are the primitives as defined by Wirth in 1983

Context switching primitives in GNU LuK

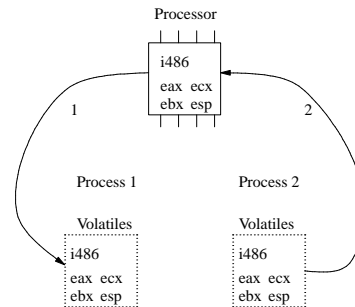
- the primitives `NEWPROCESS`, `TRANSFER` and `IOTRANSFER` are concerned with copying *Volatiles* **between process and processor**
- the procedure `TRANSFER` transfers control from one process to another process
- these primitives are *low level* primitives
 - they are normally wrapped up by higher level functions:
 - for example: `initProcess` uses `NEWPROCESS` which is similar to `new_thread` in Python

TRANSFER

- the C definition is:

```
typedef void *PROCESS;
extern void SYSTEM_TRANSFER (PROCESS *p1, PROCESS p2);
```

- and it performs the following action:



IOTRANSFER

- ```
extern void SYSTEM_IOTRANSFER (PROCESS *first,
 PROCESS *second,
 unsigned int interruptNo);
```
- the procedure `IOTRANSFER` allows process contexts to be changed when an interrupt occurs
- its function can be explained in two stages
  - firstly it transfers control from one process to another process (in exactly the same way as `TRANSFER`)
  - secondly when an interrupt occurs the processor is context switched back to the original process
- the implementation of `IOTRANSFER` involves interaction with the `FLIH`

## NEWPROCESS

- ```
extern void SYSTEM_NEWPROCESS (void (*p)(void), void *a,
                                unsigned long n,
                                PROCESS *new);
```
- `p` is a pointer to a function.
 - this function will be turned into a process
 - `a` the start address of the new processes stack
 - `n` the size in bytes of the stack
 - `new` a variable of type `PROCESS` which will contain the volatiles of the new process

How is TRANSFER implemented?

- or how do we implement a context switch?
 - first we push all registers onto the stack
 - second we need to save the current running processes stack pointer into the running process control block
 - third we need to restore the next process stack pointer into the microprocessors stack pointer
 - fourth we pop all registers from the stack

How is TRANSFER implemented?

- ```
void SYSTEM_TRANSFER (PROCESS *p1, PROCESS p2)
{
 onOrOff toOldState;

 toOldState = turnInterrupts(Off);
 asm volatile ("pusha ; pushf"); /* push all registers
 /* remember p1 is the address of a PROCESS */
 asm volatile ("movl %[p1], %%eax ; movl %%esp, (%%eax)
 :: [p1] "rm" (p1); /* p1 := top of stack
 asm volatile ("movl %[p2], %%eax ; movl %%eax, %%esp"
 :: [p2] "rm" (p2)); /* top of stack := p2
 asm volatile ("popf ; popa"); /* restore all registers
 toOldState := turnInterrupts(toOldState);
}
```

- `asm volatile`
  - means inline an assembly instruction

## How is TRANSFER implemented?

- the parameters ("movl %[p1], %%eax ; movl %%esp, (%%eax)" :: [p1] "rm" (p1));
- means
  - move p1 into register %eax
  - move %esp into the address pointed to by %eax
  - p1 is a variable which may be in a register or in memory
  - p1 is an input to the assembly instruction

## Conclusion

- we have seen the structure of a FLIHL
- we have seen how three primitives can be used to create processes, context switch between processes and react to interrupts
- we have seen how a context switch might be implemented