

Interprocess communication

- in Operating systems we find there are a number of mechanisms used for interprocess communication (IPC)
- the IPC mechanisms can be divided into two groups, those which work well using shared memory and those which work with non shared memory
- some common methods of IPC are: sockets, semaphores and mailboxes
- sockets and mailboxes are normally used by non shared memory programs
 - ie client and server on different machines

Interprocess communication in shared memory systems

- semaphores are more appropriate for multiple processes sharing some common memory
- we will be covering a semaphores and message passing after networking with sockets
- message passing
 - can be used in shared memory systems
- this week we will look at Semaphores

Semaphores: shared memory interprocess communication

- processes within an operating system do not act in isolation
 - on the one hand they co-operate to implement an application
 - on the other hand they compete for resources, processor time, device access etc
- these two elements of co-operation and competition imply some form of communication between the processes

Semaphores: shared memory interprocess communication

- in effect there are two categories for interprocess communication
 - **mutual exclusion**
 - **synchronisation**
- **mutual exclusion**
 - some resources in an operating system are non sharable, maybe access to the sound card or access to the GPU
 - access needs to be granted to one process at a time
- **synchronisation**
 - processes run asynchronously relative to each other
 - sometimes there will be points beyond which a process cannot proceed until another process has completed some activity

Mutual exclusion

- require a mechanism to ensure that only one process can manipulate data at any one time
 - *mutual exclusion*

- the concepts we discuss today are *very* important for operating systems
 - a fundamental building block

How do we implement mutual exclusion?

- simplest mechanism
 - mask processor interrupts off
 - processor cannot respond to any interrupt and therefore will execute code in sequence until it masks interrupt back on again
 - sometimes these critical sections of code are called *atomic*
 - what are this disadvantages with this approach?
 - what are this advantages with this approach?

How do we implement mutual exclusion?

- another mechanism is *semaphores*
 - essentially a binary *semaphore* is a token which can be grabbed by *only one* process at a time
 - a token is taken at the entry to the critical section and given back at the end of the critical section
 - a process can only enter once it has the token

Semaphores

- the most important single contribution towards interprocess communication was the introduction of **semaphores** by E.W. Dijkstra in 1965
 - a semaphore is a data type and the primitive operators are `wait` and `signal`

- these are the classic operators translated from Dutch words

Semaphores

- consider the following two processes:

```

/* Shared semaphore */
Semaphore token; /* initial value 1 */

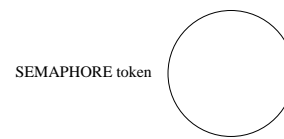
void ProcessA ()      void ProcessB ()
{
  while (TRUE) {      {
    ...               while (TRUE) {
    Wait (Token)       ...
    /* critical       Wait (Token)
    region */         /* critical
                     region */

    Signal (Token)    Signal (Token)
    ...               ...
  }                   }
}

```

Semaphores

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- Wait gets the token
- Signal returns the token

Semaphores

- note that `Wait` and `Signal` are both *atomic*
- they are implemented in software with processor interrupts masked off
- this allows us to build critical regions which can execute with processor interrupts on
- this is overall efficient as we only have to mask processor interrupts off during the execution of `Wait` and `Signal`
 - this time should be short compared with the time to execute the critical region

Semaphores

- we can express `Wait` and `Signal` in pseudo code:

```

void Wait (s)
{
  when s>0
    s--;
}

void Signal (s)
{
  s++;
}

```

Semaphores

- in our previous example the initial value of s would be 1
 - note that this is pseudo code
 - note the use of **when**

Semaphores

- we have now seen how a critical section can be achieved by using semaphore primitives `Wait` and `Signal`
- for example access to the shared buffer will be a critical section

Starting to implement a shared buffer using semaphores

- ```

void put (char ch) char get (void)
{
 Wait (Mutex) Wait (Mutex)
 /* safe to alter */ /* safe to alter */
 /* buffer */ /* buffer */
 place ch into buf remove ch from buf

 Signal (Mutex) Signal (Mutex)

 return ch;
}
 }

char buffer[Max]; /* global data */
SEMAPHORE Mutex; /* global data */

```

## Implementing synchronisation with a Semaphore

- ```

void ProcessA ()      void ProcessB ()
{
    while (TRUE) {    while (TRUE) {
        ...           ...
        Wait (sync)   /* point B */
        /* process B reached
        point B. */   Signal (sync)
        ...           ...
    }                 }
}
      
```

- we will return to this code next week

Python Semaphores and Threads

- in python you can create threads and create semaphores
 - there are a number of Python primitives which operate on semaphores but we will concentrate on those which map onto Wait and Signal

Python Semaphores and Threads

- semaphores can be created and used by:

```
from thread import start_new
from threading import Semaphore

Mutex = Semaphore(value=1)

Mutex.acquire() # Wait
Mutex.release() # Signal
```

- a thread can be created by using start_new

Example in Python of two threads synchronising

Example in Python of two threads synchronising

- `simplesync.py`

```
#!/usr/bin/env python

import sys, time
from thread import start_new
from threading import Semaphore

sync = Semaphore(value=0)

def processA (p, count):
    global sync
    print "processA", p, "comes to life"
    while True:
        time.sleep (5) # do some work
        sync.release() # indicate we have finished our w
```

- `simplesync.py`

```
def processB (p, count):
    global sync
    print "processB", p, "comes to life"
    while True:
        print "waiting for process A to complete its work"
        start_time = time.time()
        sync.acquire()
        end_time = time.time()
        print "processB", p, "spent", end_time - start_time

def main ():
    start_new(processA, (1, 0))
    processB (2, 0)

main ()
```

Example in Python of two threads implementing mutual exclusion

simplemutex.py

```
#!/usr/bin/env python

import sys, time
from thread import start_new
from threading import Semaphore

mutex = Semaphore(value=1)
n = 0 # global variable which will be incremented and
      # decremented inside the critical region
```

Example in Python of two threads implementing mutual exclusion

simplemutex.py

```
def process (p, count):
    global mutex, n
    print "process", p, "comes to life"
    while True:
        start_time = time.time()
        print "process", p, "waiting to enter"
        mutex.acquire()
        end_time = time.time()
        print "process", p, "spent", end_time - start_time, "seconds"
        # critical region
        n += 1
        if n != 1:
            print "something has gone very wrong!"
            sys.exit (1)
        time.sleep (5)
        n -= 1
        mutex.release()
        print "process", p, "finished critical region"
```

Example in Python of two threads implementing mutual exclusion

simplemutex.py

```
def main ():
    for i in range (3):
        start_new(process, (i, 0))
    process (4, 0)

main ()
```